TimeTrader: Exploiting Latency Tail to Save Datacenter Energy for Online Search

Balajee Vamanan, Hamza Bin Sohail, Jahangir Hasan, and T. N. Vijaykumar

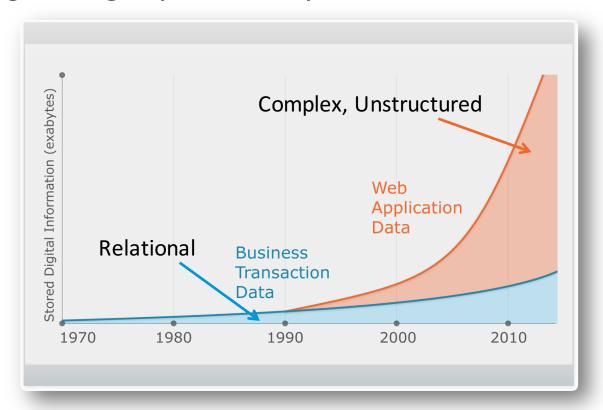






Searching large data is becoming prevalent

Data is growing exponentially



- Online Search (OLS): Interactively query and access data
 - Key component of many web applications, not only Web Search (e.g., Facebook, Twitter, advertisements)

Energy management for OLS is hard

Traditional energy management does not work

- 1. Interactive, strict service-level agreements (SLAs)
 - Cannot batch
- Short response times and inter-arrival times (e.g., 200 ms and ≈ 1 ms for Web Search)
 - Low-power modes (e.g., p-states) not applicable
- 3. Shard data over 1000s of servers \rightarrow each query searches all servers
 - Cannot consolidate workload to fewer servers

Carefully slow down (not turn off) servers without violating SLAs

Previous work

- *Pegasus* [ISCA'14]: Achieves load-proportional energy
 - Uses (global) datacenter-wide response times
 - **Shifts** latency distribution at *low* loads
- √ Saves energy at low loads
- ➤ Does not save much at high loads
 - Saves 0% at peak load
 - Datacenters operate at moderate-peak during the day (diurnal pattern) → savings desired at high loads

Pegasus saves energy at low loads but limited savings at high loads

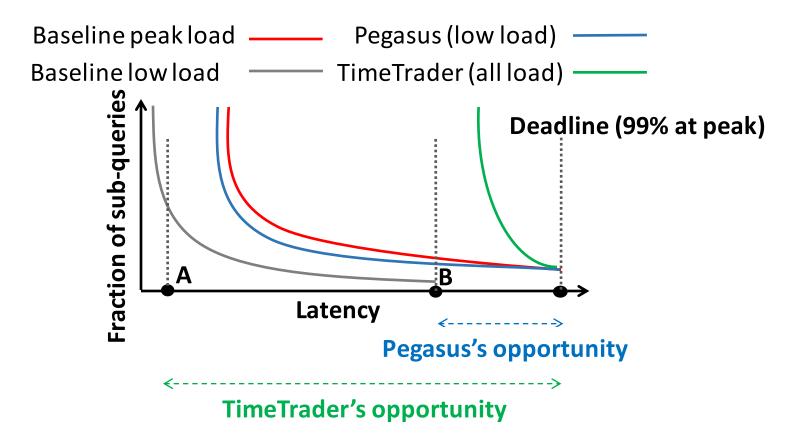
TimeTrader: contributions

Each query leads to 1000s of sub-queries Each query's budget set to accommodate 99th %-ile sub-query

- In both network and compute parts
- Key observation: 80% sub-queries complete in 20% of budget
- TimeTrader exploits both network and compute slack to save energy

TimeTrader: contributions (cont.)

Uses (local) **per-sub-query** latencies \rightarrow **reshapes** subquery response time distribution at all loads



A: median baseline at low load **B**: 99% baseline at low load

TimeTrader: contributions (cont.)

- 3. Leverages network signals to determine slack
 - Does not need fine-grained clock synchronization
- 4. Employs Earliest Deadline First (EDF) scheduling
 - Slowing down a sub-query affects other queued sub-queries
 - Critical sub-queries affected the most
 - Allows critical sub-queries to bypass others

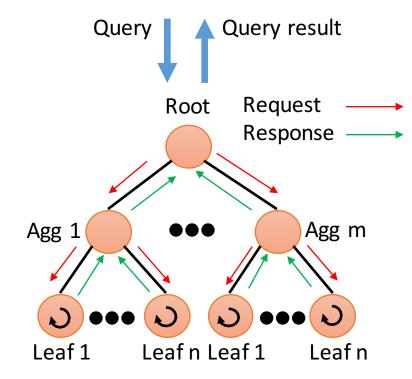
TimeTrader is the first to explore cross-layer optimization between network and architecture layers; saves 15% - 40% energy at peak - 30% loads

Outline

- Introduction
- Background
 - OLS architecture
 - Tail latencies and SLA budgets
- TimeTrader's design
- Methodology
- Results

Background: OLS architecture

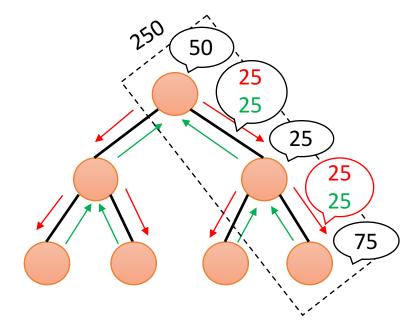
- Partition-aggregate (tree)
- Request-compute-reply
 - Request: root → leaf
 - Compute: leaf node
 - Reply: leaf → root



- Root waits until deadline to generate overall response
 - Some replies may take longer in network or compute
 - Collisions in network → TCP retransmit
 - Imperfect sharding → disproportionately large compute
 - Long tails → tail latency problem

OLS tail latencies and SLA budgets

- Missed replies due to tail latency affect quality
 - e.g., SLA of 1% missed deadlines
 - Time budgets based on 99th %-ile latency of sub-queries
- To optimize network and compute separately
 - Split budget into network and compute
 - e.g., total = 250 ms→ flow deadlines 25 ms



80% of sub-queries complete within 20% of budget → slack

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- Introduction
- Background
- TimeTrader's design
 - Key ideas
 - Determining slack
 - Slowing down based on slack
 - EDF
- Methodology
- Results

TimeTrader: design

- TimeTrader exploits per-sub-query slack
 - Sources of slack
 - Network
 - ✓ Request
 - Reply (after compute, unpredictable)
 - Compute
 - Actual compute (query dependent)
 - ✓ Queuing (due to local load variations)
- 1. Determine (a) request slack and (b) compute slack
- Slow down based on slack and load
- 3. Shield critical sub-queries from slowed down sub-queries
 - EDF implementation
- Intel's Running Avg. Power Limit (RAPL) to set power states

1(a): Determine request slack

- Requests traverse network from root to leaf
- Timestamp at parent and leaf?
 - **×**Clock skew ≈ slack
 - ✓ Estimate based on already-available network signals
 - e.g., Explicit Congestion Notification (ECN) marks,TCP timeouts

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If a leaf doesn't see either ECN or Timeouts

*Request Slack = Network Budget - Median Latency

Else

*Request Slack = 0
```

1(b): Determine compute-queuing slack

- Exploit local load variations in each leaf server
- Slack calculation
 - Details in the paper

Total Slack = Request Slack + Compute-Queuing Slack

2: Slow down based on slack

- Given the slack, calculate slow down factor
 - As a fraction of compute budget
- But not all of the slack can be used
 - Slack calculation assumes no queuing
 - Need to scale slack based on load

Slowdown factor = Total Slack * Scale / Compute Budget

- Scale depends on load
 - Feedback controller (details in the paper)

3: Shield critical sub-queries

- Earliest Deadline First (EDF)
 - Implementation details in the paper

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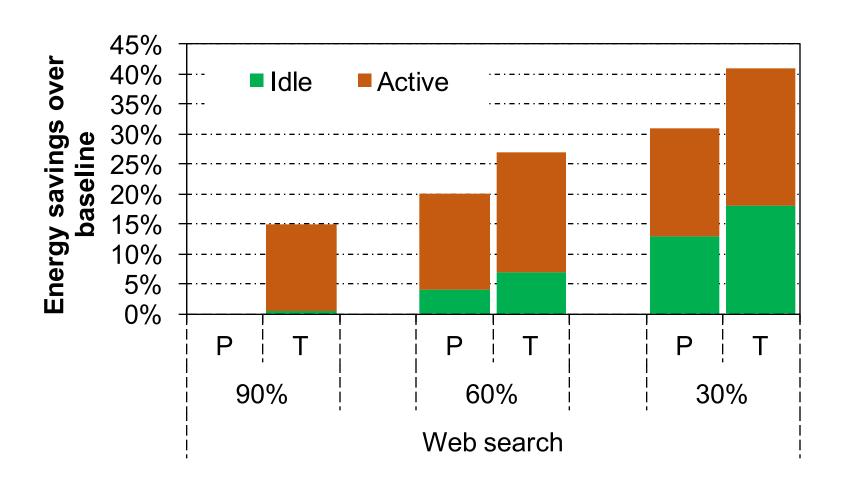
Methodology

- Compute
 - Real measurements for service time, power
- Network
 - Real measurements at a small scale
 - Tails effects only in large clusters → simulate with ns-3

Workload: Web Search (Search) from CloudSuite 2.0

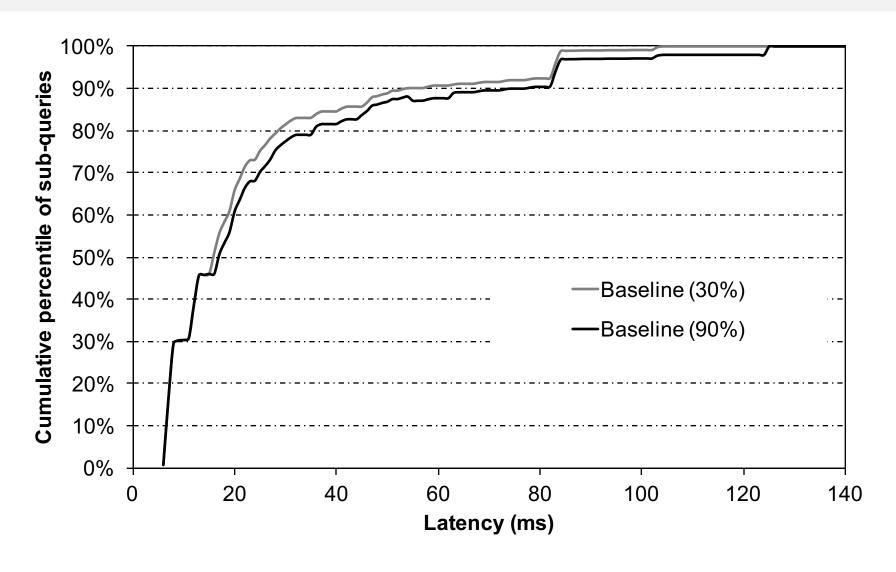
- Search index from Wikipedia
- 3000 queries/s at peak load → 90% load
- Deadline budget: Overall 125 ms
 - Network (request/reply): 25 ms
 - Compute: 75 ms
- SLA of 1% missed deadlines

Idle and Active Energy Savings

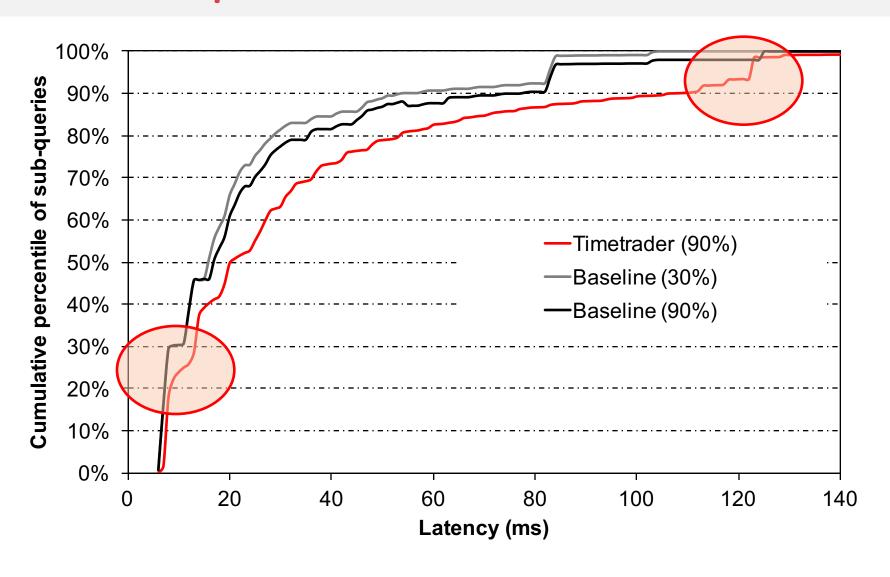


TimeTrader saves 15% - 40% energy over baseline (17% over Pegasus at 30% load)

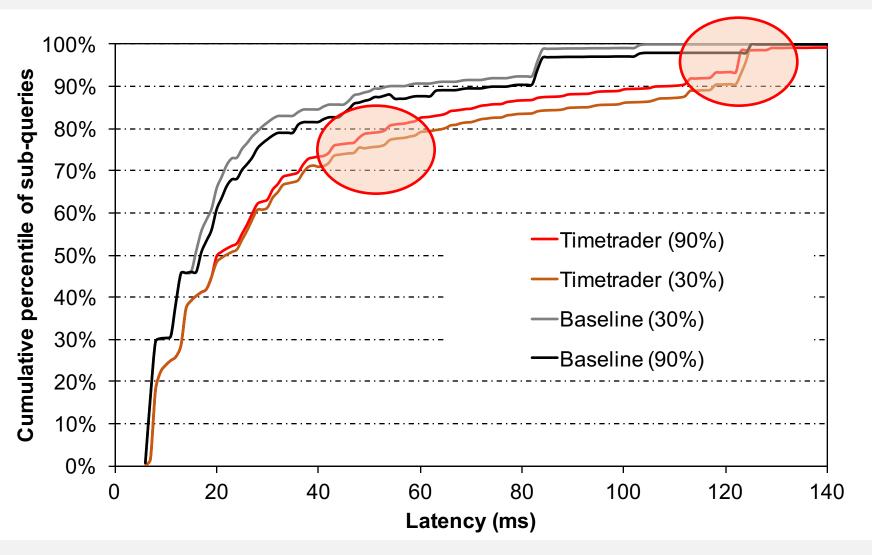
Response time distribution



Response time distribution



Response time distribution



TimeTrader reshapes distribution at all loads; Slows ~80% of requests at all loads

Conclusion

TimeTrader...

- Exploits sub-query slack
- Reshapes response time distribution at all loads
 - Saves 15% energy at peak, 40% energy at 30% load
- Leverages network signals to estimate slack
- Employs EDF to decouple critical sub-queries from subcritical sub-queries

TimeTrader converts the performance disadvantage of latency tail into an energy advantage!