

HETEROGENEOUS SYSTEM COHERENCE FOR INTEGRATED CPU-GPU SYSTEMS

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Hardware coherence can increase the utility of heterogeneous systems

Major bottlenecks in current coherence implementations

- High bandwidth difficult to support at directory
- -Extreme resource requirements

We propose Heterogeneous System Coherence

- -Leverages spatial locality and region coherence
- Reduces bandwidth by 94%
- -Reduces resource requirements by 95%

















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LOGICAL INTEGRATION



General-purpose GPU computing

- -OpenCL
- -CUDA

Heterogeneous Uniform Memory Access (hUMA) –Shared virtual address space –Cache coherence

▲ Allows new heterogeneous apps



OUTLINE

Motivation

Background

- -System overview
- -Cache architecture reminder
- Heterogeneous System Bottlenecks
- Heterogeneous System Coherence Details
- Results
- Conclusions



SYSTEM LEVEL



SYSTEM OVERVIEW





SYSTEM OVERVIEW





SYSTEM OVERVIEW





CACHE ARCHITECTURE REMINDER CPU/GPU L2 CACHE





DIRECTORY ARCHITECTURE REMINDER





BACKGROUND SUMMARY

▲ System under investigation

- -Heterogeneous CPU-GPU on chip
- -High-bandwidth DRAM

Directory pipeline complex

- -MSHR array is associative
- -Difficult to pipeline with more than 1 request per cycle
- -Important resources: MSHR entries



OUTLINE

Motivation

Background

Heterogeneous System Bottlenecks

- -Simulation overview
- Directory bandwidth
- –MSHRs
- -Performance is significantly affected
- Heterogeneous System Coherence Details
- Results

Conclusions

SIMULATION DETAILS

▲ gem5 simulator

- -Simple CPU
- -GPU simulator based on AMD GCN
- -All memory requests through gem5

Workloads

- -Modified to use hUMA
- -Rodinia & AMD APP SDK

CPU Clock	2 GHz
CPU Cores	2
CPU Shared L2	2 MB (16-way banked)
GPU Clock	1 GHz
Compute Units	32
GPU Shared L2	4 MB (64-way banked)
L3 (Memory-side)	16 MB (16-way banked)
DRAM	DDR3, 16 channels
Peak Bandwidth	700 GB/s
Baseline Directory	256k entries (8-way banked)



GPGPU BENCHMARKS

Rodinia benchmarks

- **bp** trains the connection weights on a neural network
- bfs breadth-first search
- hs performs a transient 2D thermal simulation (5-point stencil)
- lud matrix decomposition
- nw performs a global optimization for DNA sequence alignment
- km does k-means clustering
- sd speckle-reducing anisotropic diffusion

AMD SDK

- **bn** bitonic sort
- dct discrete cosine transform
- **hg** histogram
- mm matrix multiplication



SYSTEM BOTTLENECKS



DIRECTORY TRAFFIC





RESOURCE USAGE



PERFORMANCE OF BASELINE

COMPARED TO UNCONSTRAINED RESOURCES

BOTTLENECKS SUMMARY

Directory bandwidth

- -Must support up to 4 requests per cycle
- -Difficult to construct pipeline

Resource usage

- -MSHRs are a constraining resource
- -Need more than 10,000
- –Without resource constraints, up to 4x better performance

OUTLINE

Motivation

Background

Heterogeneous System Bottlenecks

Heterogeneous System Coherence Details

- Overall system design
- Region buffer design
- Region directory design
- -Example
- -Hardware complexity
- Results

Conclusions

BASELINE DIRECTORY COHERENCE

HETEROGENEOUS SYSTEM COHERENCE (HSC) AMD

HETEROGENEOUS SYSTEM COHERENCE (HSC) AMD

HSC: EXAMPLE MEMORY REQUEST

HSC: L2 CACHE & REGION BUFFER

HSC: REGION DIRECTORY

HSC: HARDWARE COMPLEXITY

Region protocols reduce directory size

-Region directory: 8x fewer entries

Region buffers

- -At each L2 cache
- -1-KB region (16 64-B blocks)
- –16-K region entries
- Overprovisioned for low-locality workloads

HSC SUMMARY

Key insight

- -GPU-CPU applications exhibit high spatial locality
- -Use direct-access bus present in systems
- -Offload bandwidth onto direct-access bus
- ▲ Use coherence network only for permission
- Add region buffer to track region information
 - –At each L2 cache
 - -Bypass coherence network and directory
- Replace directory with region directory
 - -Significantly reduces total size needed

OUTLINE

Motivation

- Background
- Heterogeneous System Bottlenecks
- Heterogeneous System Coherence Details

Results

- -Speed-up
- -Latency of loads
- –Bandwidth
- –MSHR usage
- Conclusions

THREE CACHE-COHERENCE PROTOCOLS

Broadcast: Null-directory that broadcasts on all requests

Baseline: Block-based, mostly inclusive, directory

HSC: Region-based directory with 1-KB region size

HSC PERFORMANCE

DIRECTORY TRAFFIC REDUCTION

HSC RESOURCE USAGE

▲ Used a detailed timing simulator for CPU and GPU

▲ HSC significantly improves performance

- -Reduces the average load latency
- Decreases bandwidth requirement of directory

▲ HSC reduces the required MSHRs at the directory

RELATED WORK

Coarse-grained coherence

- -Region coherence
 - Applied to snooping systems [Cantin, ISCA 2005] [Moshovos, ISCA 2005]
 [Zebchuk, MICRO 2007]
 - Extended to directories [Fang, PACT 2013] [Zebchuk, MICRO 2013]
- -Spatiotemporal coherence [Alisafaee, MICRO 2012]
- -Dual-grain directory coherence [Basu, UW-TR 2013]
 - Primarily focused on directory size
- GPU coherence [Singh et al. HPCA 2013]
 - -Intra-GPU coherence

CONCLUSIONS

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Questions?

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